A stronger square conjecture on binary words

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Square packing What is square?

Square

A square is a word (sequence of letters) of the form xx.

Example

Squares	Non-squares
аа	aba
abab	bbab
baabbaab	baabbaaa
abaaababaaab	ааааааа
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Square packing Squares on a word

Counting rule

Don't count the same square twice or more times!

What squares does the following word (length 16) contain?

abbaaaabaaaabaa

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• aa (occurs 8 times but count is just 1)

- A - E - M

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- 4 E 6 4 E 6

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abba<u>aaabaaaabaa</u>

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Square packing Squares on a word

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- aaaabaaaaba
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- bb
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The count is 6

4 E 6 4 E

Square packing Counting squares: example

Exercise

How many squares does the following word (length 20) contain?

Square packing Counting squares: example

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Answer

10 $(a^2, a^4, a^6, a^8, a^{10}, a^{12}, a^{14}, a^{16}, a^{18}, a^{20}).$

Square packing Counting squares: example

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Answer

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$$(a^2, a^4, a^6, a^8, a^{10}, a^{12}, a^{14}, a^{16}, a^{18}, a^{20}).$$

Proposition

The unary word of length *n*, that is, a^n , contains exactly $\lfloor n/2 \rfloor$ squares.

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Square packing Counting squares: example

Exercise

How many squares does the following word (length 20) contain?

abaabaaabaabaaabaaaa

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Square packing Counting squares: example

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Answer

13 $(a^2, (aa)^2, (aab)^2, (aba)^2, (baa)^2, (aaba)^2, (abaa)^2, (baaa)^2, (aabaaab)^2, (abaaabaa)^2, (abaaabaa)^2, (baaabaa)^2, (baaabaaa)^2).$

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Square packing Counting squares: example

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Answer

Thus, binary words outperform unary words in packing squares.

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Square packing Quiz-setting

Square packing

Find a word that contains many squares relative to its length.

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Upper bounds Linear bounds; proved and conjectured

Theorem ([Fraenkel & Simpson 1998])

A word of length n contains at most 2n squares.

Idea of simple proof [Ilie 2005]

Ignore all but the rightmost occurrence of each square, and we can see that, at each position, at most 2 rightmost occurrences can start. $\hfill \Box$

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Square conjecture

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Upper bounds Conjecture

Unsuccessful attempt

If we could show that "at each position of a word, at most one rightmost occurrence could start," then the conjecture would be proved.

Counterexample

Consider *abaababaab*. The rightmost occurrences of *abaaba* and *abaababaab* start both on position 1.

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Square-dense words Examples

Known square-dense words include

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Square-dense words Examples

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Square-dense words Examples

Known square-dense words include

- the following word by Fraenkel and Simpson:
- Let us propose a "simpler" one as:

 $w_{\rm jms} = abaabaaabaaaabaaaabaaaab \cdots = a^1 ba^2 ba^3 ba^4 ba^5 b \cdots$

Square-dense words Examples

Known square-dense words include

- the following word by Fraenkel and Simpson:
- Let us propose a "simpler" one as:

 $w_{\rm jms} = abaabaaabaaaabaaaabaaaab \cdots = a^1 ba^2 ba^3 ba^4 ba^5 b \cdots$

Note that both of these words are binary.

Square-densest words Candidates are binary

Conjecture

Square-densest words are binary.

binary words with at most $1 \ b$ is square-sparse

Recall that the word a^n contains exactly $\lfloor n/2 \rfloor$ squares. Consider a word $a^i b a^{n-i-1}$ with 1 *b*. The *b* cannot make contribution to any square, so replacing it with *a* does not destroy any square. Hence, it contains at most $\lfloor n/2 \rfloor$ squares.

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Stronger conjecture Stronger conjecture on binary words

Let us strengthen the square conjecture by the # of occurrences of b's as follows.

Stronger conjecture on binary words

For $k \ge 2$, a binary word w of length n with k b's contains at most $\frac{2k-1}{2k+2}n$ squares.

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Stronger conjecture Asymptotic tightness

Theorem

Both of the words $w_{\rm fs}$ and $w_{\rm jms}$ achieve the tighter bound asymptotically.

Inductive approach to the stronger conjecture Preliminaries

A binary word w with k b's can be represented as

$$w=a^{i_1}ba^{i_2}b\cdots a^{i_k}ba^{i_{k+1}}$$

for some coefficients $i_1, \ldots, i_{k+1} \ge 0$.

Coefficient sequence

The coefficient sequence of w means the sequence i_1, \ldots, i_{k+1} .

Coefficient set

The *coefficient set* of w, denoted by I(w), is the **multiset** $\{i_1, i_2, \ldots, i_k, i_{k+1}\}$.

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Inductive approach to the stronger conjecture sufficient condition inequality

Let

$$I(w)[j]$$
 the *j*-th smallest element of $I(w)$
 $I(w)[max]$ the maximum element of $I(w)$
 $\#Sq(w)$ the number of squares on the word w

Lemma

If a word w_k of length n with k b's satisfies

$$\#\operatorname{Sq}(w_k) \leq \left\lfloor \frac{I(w_k)[\max]}{2} \right\rfloor + \sum_{j=1}^{|I(w_k)|-2} (I(w_k)[j]+1),$$

then #Sq $(w_k) \leq \frac{2k-1}{2k+2}n$.

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Inductive approach to the stronger conjecture Induction strategy

Induction hypothesis

The inequality holds for all words with at most k - 1 b's.

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The inequality holds for all words with at most k - 1 b's.

• Apply an operation to a word w_{k-1} with k-1 b's to yield a word w_k with k b's.

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Inductive approach to the stronger conjecture Induction strategy

Induction hypothesis

The inequality holds for all words with at most k - 1 b's.

- Apply an operation to a word w_{k-1} with k-1 b's to yield a word w_k with k b's.
- Prove that the # of squares thus created is small enough to keep the inequality valid.

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Inductive approach to the stronger conjecture Induction strategy: example

Induction hypothesis

 $w_{k-1} = a^1 b a^2 b \cdots a^{k-1} b a^k$ (with k-1 b's) satisfies the inequality, that is, $\# \operatorname{Sq}(w_{k-1}) \leq \lfloor k/2 \rfloor + \sum_{j=1}^{k-2} (j+1)$.

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$$w_{k-1} \leftarrow ba^{k+1} \Rightarrow w_k = a^1 ba^2 b \cdots a^{k-1} ba^k ba^k a.$$

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$$w_{k-1} \leftarrow ba^{k+1} \Rightarrow w_k = a^1 ba^2 b \cdots a^{k-1} ba^k ba^k a.$$

2 The catenation creates
$$(k - 1) + 1$$
 squares, that is,
 $(a^{k-1}ba)^2, (a^{k-2}ba^2)^2, \dots, (aba^{k-1})^2, (ba^k)^2.$

Inductive approach to the stronger conjecture Induction strategy: example

Induction hypothesis

 $w_{k-1} = a^1 b a^2 b \cdots a^{k-1} b a^k$ (with k-1 b's) satisfies the inequality, that is, $\# \operatorname{Sq}(w_{k-1}) \leq \lfloor k/2 \rfloor + \sum_{j=1}^{k-2} (j+1)$.

$$w_{k-1} \leftarrow ba^{k+1} \Rightarrow w_k = a^1 ba^2 b \cdots a^{k-1} ba^k ba^k a.$$

• The catenation creates (k-1)+1 squares, that is, $(a^{k-1}ba)^2, (a^{k-2}ba^2)^2, \dots, (aba^{k-1})^2, (ba^k)^2$.

In Thus,

$$\begin{split} \# \mathrm{Sq}(w_k) &\leq \quad \lfloor (k+1)/2 \rfloor + \sum_{j=1}^{k-2} (j+1) + (k-1) + 1 \\ &= \quad \lfloor (k+1)/2 \rfloor + \sum_{j=1}^{k-1} (j+1) \end{split}$$

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Inductive approach to the stronger conjecture Theorems

1) When the coefficients are pairwise-distinct

Theorem

For $k \ge 2$, let $w_k = a^{i_0} b a^{i_1} b \cdots b a^{i_k}$ be a binary word of length n with k b's. If the coefficients i_1, \ldots, i_{k-1} are pairwise-distinct, then #Sq $(w_k) \le \frac{2k-1}{2k+2}n$.

Recall $w_{jms} = a^1 b a^2 b a^3 b a^4 b \dots$ Thus, the tighter bound holds for any subword of w_{jms} .

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Inductive approach to the stronger conjecture Theorems

2) When the coefficient sequence is sorted

Theorem

For $k \ge 2$, let $w_k = a^{i_0} b a^{i_1} b \cdots b a^{i_{k-1}} b a^{i_k}$ be a binary word of length n with k b's. If $i_1 \le i_2 \le \cdots \le i_{k-1}$, then $\#\operatorname{Sq}(w_k) \le \frac{2k-1}{2k+2}n$.

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Inductive approach to the stronger conjecture Theorems

3) Pairwise-distinct extrema

A maximum of the coefficient sequence i_0, \ldots, i_k is its subsequence $i_{\ell}, i_{\ell+1}, \ldots, i_{r-1}, i_r$ such that $i_{\ell} < i_{\ell+1} = \cdots = i_{r-1} > i_r$. The notion of minima is defined analogously.

Theorem

For $k \ge 2$, let w_k be a binary word of length n with k b's. If all maxima of its coefficient sequence are pairwise-distinct, then $\#\operatorname{Sq}(w_k) \le \frac{2k-1}{2k+2}n$.

Recall $w_{\rm fs} = a^1 ba^2 ba^3 ba^2 ba^3 ba^4 ba^3 ba^4 ba^5 b \cdots$, and its maxima $i_2 i_3 i_2, i_3 i_4 i_3, i_4 i_5 i_4, \ldots$ are pairwise-distinct. Hence, the tighter bound holds for any subword of $w_{\rm fs}$.

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Inductive approach to the stronger conjecture Theorems

4) Small # of *b*'s

proposition

For $2 \le k \le 9$, any word *w* of length *n* with at most *k b*'s satisfies #Sq(*w*) $\le \frac{2k-1}{2k+2}n$.

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Summary of the solved cases

The tighter bound holds if one of the following conditions is safisfied:

- Extrema of the coefficient sequence are pairwise-distinct, including the cases when:
 - the coefficients are pairwise-distinct
 - the sequence is sorted
- Ø Multiplicity of coefficients is at most 4.

$$|w|_b \leq 9.$$

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Open problems

- Prove the bound $\frac{2k-1}{2k+2}n$ in general,
- or Find a counterexample (a binary word of length n with k b's that contains more than ^{2k-1}/_{2k+2} n squares).
- Prove that square-densest words are binary.
- Generalize the bound for arbitrary alphabets.

Acknowledgements

Thank you very much for your attention!

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